

# Time-parallelization of sequential data assimilation problems

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Luenberger observer

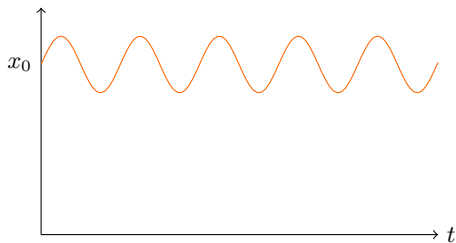
Time-parallelization setting

Parareal algorithm

Diamond strategy (Parareal case)

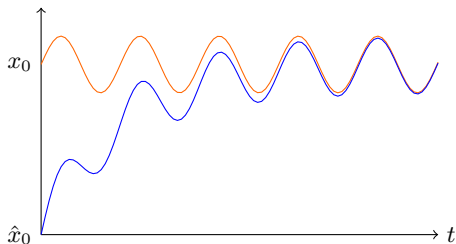
A textbook example

$$(1) \quad \begin{cases} \dot{x}(t) = Ax(t) + Bu(t), \\ x(0) = x_0 \text{ unknown}, \\ y(t) = Cx(t). \end{cases}$$



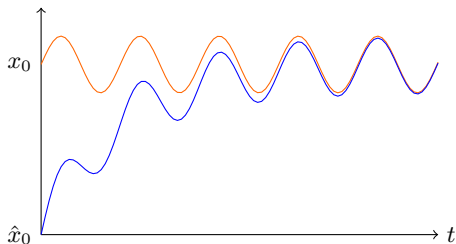
$$(1) \quad \begin{cases} \dot{x}(t) = Ax(t) + Bu(t), \\ x(0) = x_0 \text{ unknown}, \\ y(t) = Cx(t). \end{cases}$$

$$(2) \quad \begin{cases} \dot{\hat{x}}(t) = A\hat{x}(t) + Bu(t) \\ \quad \quad \quad + L[y(t) - \hat{y}(t)], \\ \hat{x}(0) = \hat{x}_0, \\ \hat{y}(t) = C\hat{x}(t). \end{cases}$$



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Is there a *convenient* way to choose the *observer gain*  $L$  ?

Note that

$$(2) \iff \begin{cases} \dot{\hat{x}}(t) = (A - LC)\hat{x}(t) + (Bu(t) + Ly(t)), \\ \hat{x}(0) = \hat{x}_0. \end{cases}$$

and then  $x(t) - \hat{x}(t) = e^{(A-LC)t} (x(0) - \hat{x}(0))$

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## THEOREM ► IDENTITY OBSERVER THEOREM [LUENBERGER]

Given a completely *observable* system (1), an identity observer of the form (2) can be constructed, and the coefficients of the **characteristic polynomial of the observer** can be selected arbitrarily.

Note that

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## PROPOSITION 1.1

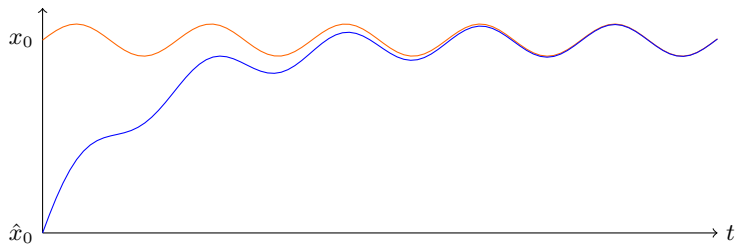
We assume System (1) is observable and the eigenvalues of  $A - LC$  are negative and simple. Then, we have

$$\|e^{(A-LC)t}\| \leq \gamma e^{-\mu t}$$

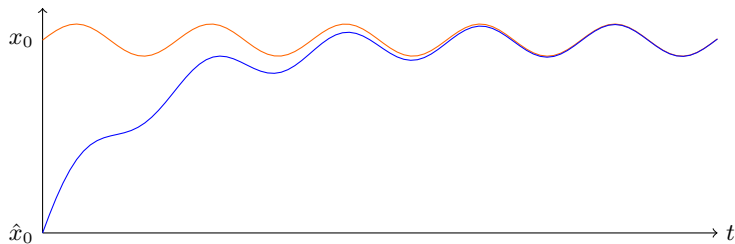
with  $\mu := \min_{\nu \in \sigma(A-LC)} |\nu|$  and  $\gamma := \text{cond}(V) = \|V^{-1}\| \|V\|$ , where  $V$  is the matrix whose rows are the eigenvectors of  $A - LC$  and  $\|\cdot\|$  represents the induced 2-norm of a matrix.



# DIAMOND STRATEGY



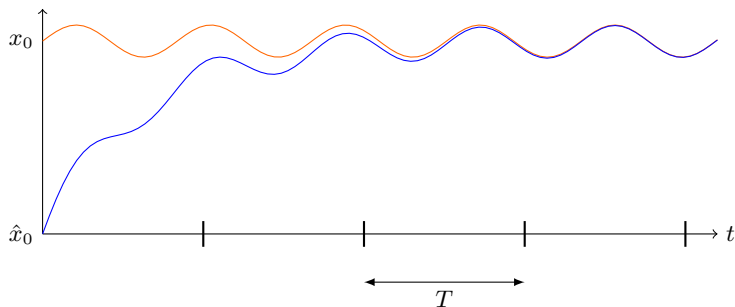
# DIAMOND STRATEGY



Luenberger observer

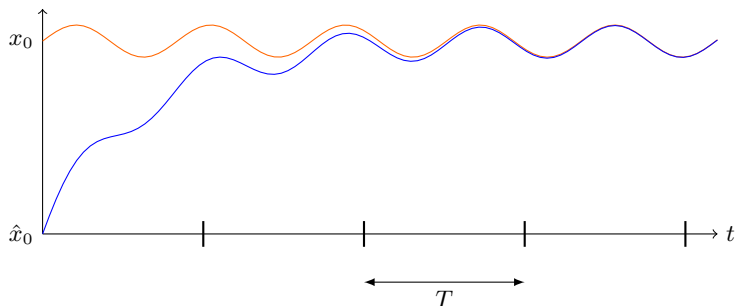
Time-parallel method

# DIAMOND STRATEGY



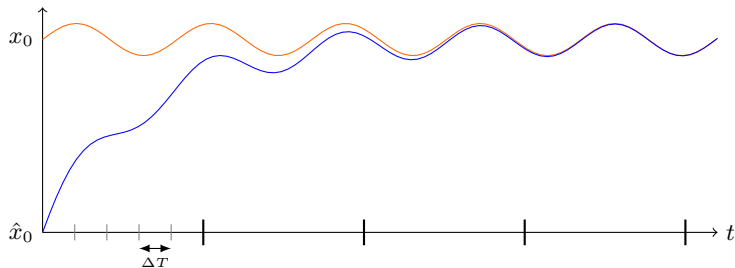
- Divide the time interval into *windows*  $W_\ell$  of a given length  $T > 0$ .

# DIAMOND STRATEGY



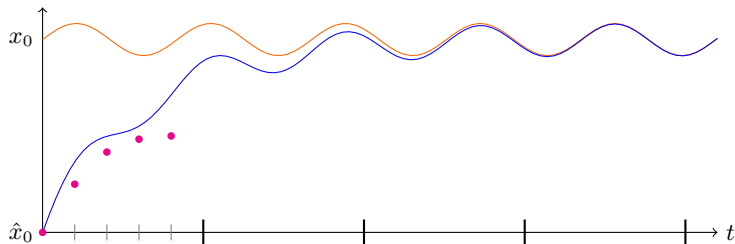
- ▶ Divide the time interval into *windows*  $W_\ell$  of a given length  $T > 0$ .
- ▶ Solve Equation (2) on each window, in a sequential order, using a time-parallel algorithm.

## DIAMOND STRATEGY (TIME-PARALLELIZATION)



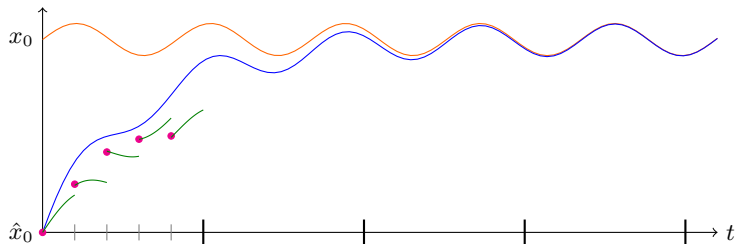
- Decompose  $W_\ell$  into  $N$  *subintervals* of length  $\Delta T$ .

## DIAMOND STRATEGY (TIME-PARALLELIZATION)



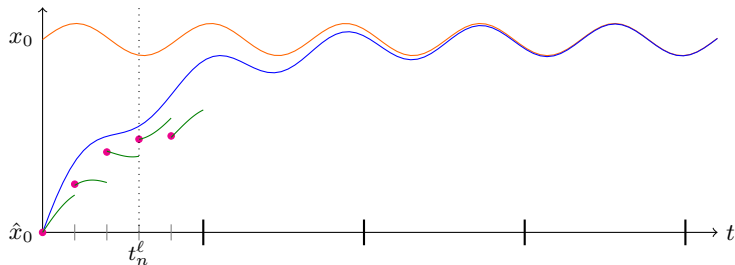
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## DIAMOND STRATEGY (TIME-PARALLELIZATION)



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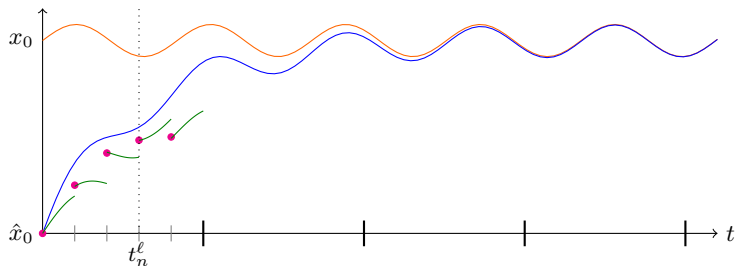


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$$J_{\ell,n}^h := \hat{X}_{\ell,n}^h - \hat{x}_{\parallel}(t_n^{\ell-}).$$



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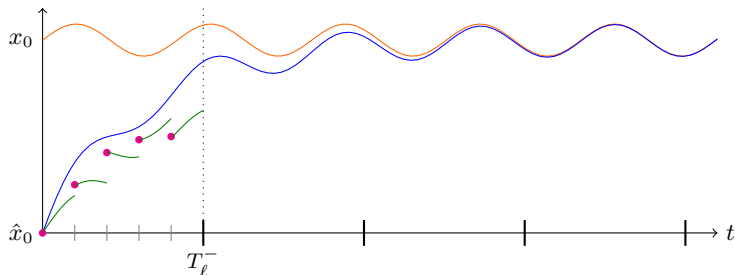


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Next step : define a suitable stopping criterion !

# DIAMOND STRATEGY (STOPPING CRITERION)



## LEMMA

Under the assumptions of Proposition 1.1, we have

$$\|(x - \hat{x}_{\parallel})(T_{\ell}^{-})\| \leq \gamma \left( \|x_0 - \hat{x}_0\| + \sum_{i=1}^{\ell} e^{-\mu(i-1)T} \cdot \gamma \sum_{n=1}^{N-1} e^{\mu n T} \|J_{i,n}^h\| \right) e^{-\mu \ell T}$$

## PROPOSITION 2.1 ► A POSTERIORI ESTIMATE

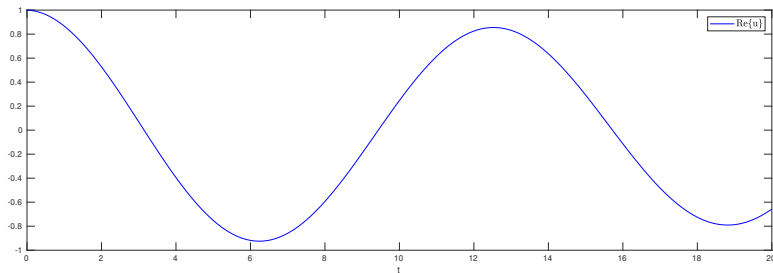
Let us assume that  $h$  is obtained from the stopping criterion in  $W_\ell$

$$2\gamma \sum_{n=1}^{N-1} e^{\mu n \Delta T} \|J_{\ell,n}^h\| \leq \tilde{\gamma} \frac{e^{-\mu(\ell-1)T}}{2^{\ell-1}}$$

where  $\tilde{\gamma}$  is an arbitrary parameter. Then, the rate of convergence of  $\hat{x}_\parallel(t)$  to  $x(t)$  is bounded by  $\mu$ , i.e.

$$\|(x - \hat{x}_\parallel)(T_\ell^-)\| \leq \gamma (\|x(0) - \hat{x}(0)\| + \tilde{\gamma}) e^{-\mu \ell T}.$$

# PARAREAL ALGORITHM



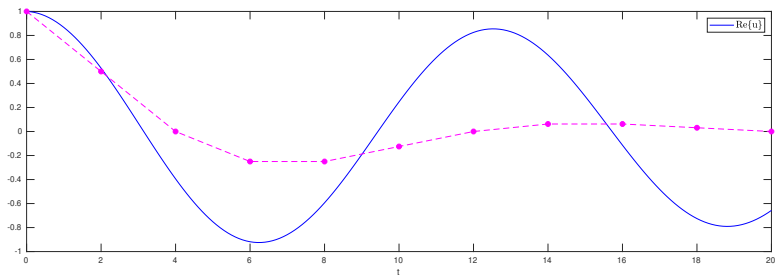
Dahlquist equation  $\dot{u}(t) = -\frac{i}{2}u$  in  $[0, 20]$

To solve the problem

$$\begin{cases} \dot{u}(t) = f(u(t)), & t \in [0, T] \\ u(0) = u_0 \end{cases}$$

we decompose the time interval on  $N$  subintervals, denoted by  $(t_{n-1}, t_n)$ .

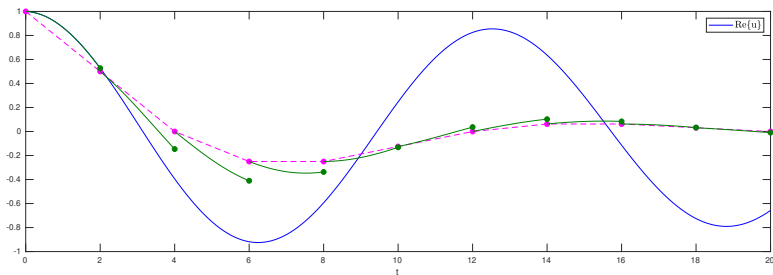
# PARAREAL ALGORITHM



- Impose arbitrary values on the subintervals by using the **coarse solver**  $\mathcal{G}$ :

$$U_0^0 = u_0, U_n^0 = \mathcal{G}(t_n, t_{n-1}, U_{n-1}^0).$$

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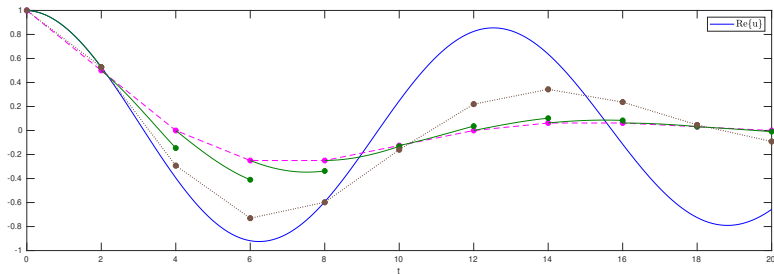
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- Using the **fine solver**  $\mathcal{F}$ , solve in parallel

$$\begin{cases} \dot{u}(t) = f(u(t)), & t \in [t_{n-1}, t_n] \\ u(t_{n-1}) = U_{n-1}^0. \end{cases}$$

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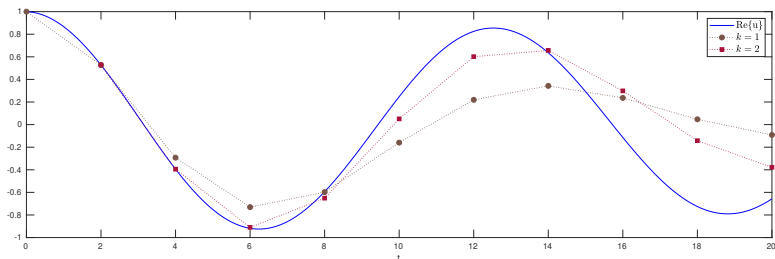
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- Smooth the discontinuities previously introduced by defining

$$U_n^1 := \mathcal{F}(t_n, t_{n-1}, U_{n-1}^0) + \mathcal{G}(t_n, t_{n-1}, U_{n-1}^1) - \mathcal{G}(t_n, t_{n-1}, U_{n-1}^0).$$

# PARAREAL ALGORITHM



At iteration  $k$ :

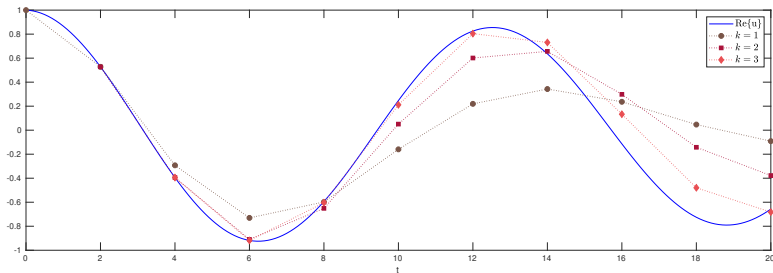
- compute  $\{\mathcal{F}(t_n, t_{n-1}, U_{n-1}^{k-1})\}_{n=1}^N$  in parallel.
- Update the sequence

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What about its convergence ?

THEOREM ► CONVERGENCE OF PARAREAL [GANDER AND HAIRER]

(...) at iteration  $k$  of the Parareal algorithm, we have the bound

$$\|u(t_n) - U_n^k\| \leq \frac{C_3}{C_1} \frac{(C_1 \Delta T^{p+1})^{k+1}}{k!} (1 + C_2 \Delta T)^{n-(k+1)} \prod_{j=0}^k (n - j).$$

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► **Superlinear** rate of convergence.

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- Superlinear rate of convergence.
- Among other assumptions,  $\mathcal{F}(t_n, t_{n-1}, U_{n-1}^k)$  is the exact solution on  $(t_{n-1}, t_n)$ , and  $\mathcal{G}$  must satisfy

$$\|\mathcal{G}(t + \Delta T, t, x) - \mathcal{G}(t + \Delta T, t, y)\| \leq (1 + C_2 \Delta T) \|x - y\|.$$

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- The result is well suited for **non-decaying problems**.

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## THEOREM 3.1 ► CONVERGENCE OF PARAREAL FOR DECAYING PROBLEMS

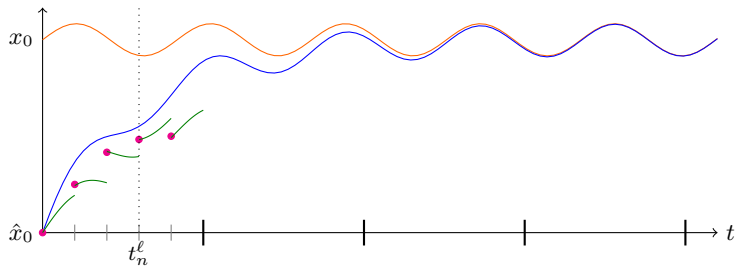
(...) We also assume that  $\mathcal{F}$  and  $\mathcal{G}$  are Lipschitz with respect to the initial conditions:

$$\max \{ \|\mathcal{F}(t_n, t_{n-1}, y) - \mathcal{F}(t_n, t_{n-1}, z)\|, \|\mathcal{G}(t_n, t_{n-1}, y) - \mathcal{G}(t_n, t_{n-1}, z)\| \} \leq \varepsilon \|y - z\|,$$

for a constant  $\varepsilon \in (0, 1)$ . Then, after  $k$  iterations of the Parareal algorithm, we have

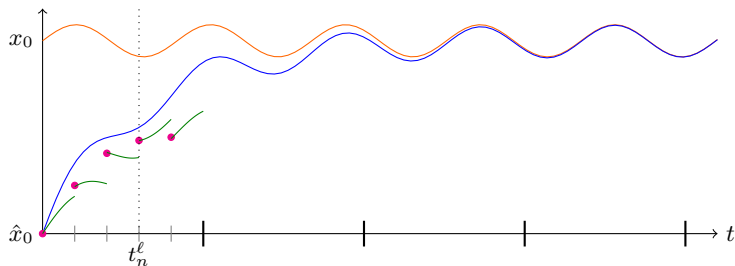
$$\|U_n^k - u(t_n)\| \leq \begin{cases} 0 & n \leq k \\ B_n^k := \alpha \beta^k \sum_{i=0}^{n-k-1} \binom{k+i}{k} \varepsilon^i & n > k. \end{cases}$$

## DIAMOND STRATEGY (PARAREAL CASE)



- The Luenberger observer  $\hat{x}(t)$  is a decaying problem (Proposition 1.1).

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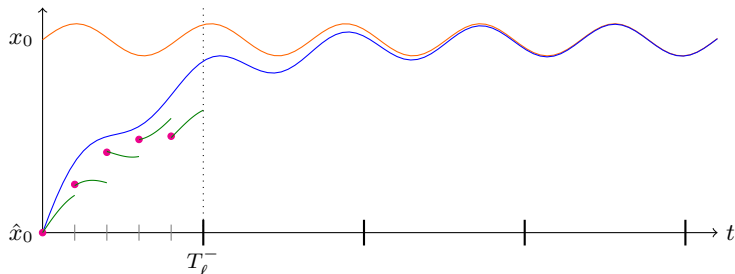


- The Luenberger observer  $\hat{x}(t)$  is a decaying problem (Proposition 1.1).
- The number of parareal iterations  $\{k_\ell\}_\ell$  can be determined from
  - (a) Proposition 2.1 (*a posteriori* estimate)

$$2\gamma \sum_{n=1}^{N-1} e^{\mu n \Delta T} \left\| \hat{X}_{\ell,n}^{k_\ell} - \hat{x}_\parallel(t_n^\ell) \right\| \leq \tilde{\gamma} \frac{e^{-\mu(\ell-1)T}}{2^{\ell-1}}, \quad \tilde{\gamma} \text{ arbitrary.}$$



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- (b) Theorem 1.3 (*a priori* bound)

$$\left\| \hat{X}_{\ell,n}^{k_\ell} - \hat{x}_\parallel(t_n^\ell) \right\| \leq B_n^k(\alpha, \beta, \varepsilon)$$

# DIAMOND STRATEGY (PARAREAL CASE)

## THEOREM 4.1

We keep the assumptions of Proposition 1.1 and Theorem 1.3. For a window  $W_\ell$  and  $\tilde{\gamma} > 0$ , we define

$$k_\ell = \begin{cases} \min S_\ell & S_\ell \neq \emptyset \\ k_{\ell-1} & S_\ell = \emptyset \end{cases}$$

where

$$S_\ell = \left\{ k \in \mathbb{N}^*, k \leq N-1 : 2\gamma \sum_{n=1}^{N-1} e^{\mu n \Delta T} B_n^k(\alpha, \beta, \varepsilon) \leq \tilde{\gamma} \frac{e^{-\mu(\ell-1)T}}{2^{\ell-1}} \right\}.$$

Suppose that we apply the *Diamond strategy* using  $k_\ell$  iterations of the Parareal algorithm. Then, the stopping criterion is satisfied.

Consider the second order system

$$\begin{cases} \ddot{z} + 2\eta\omega\dot{z} + \omega^2 z = 5 + e^{-100t} \sin(\frac{3}{4}t), \\ z(0), \dot{z}(0) \text{ unknown,} \end{cases}$$

with  $\dot{z}(t)$  as output.

Consider the second order system (in matricial form)

$$\begin{cases} \dot{x} = \begin{bmatrix} 0 & 1 \\ -\omega^2 & -2\eta\omega \end{bmatrix} x + \begin{pmatrix} 0 \\ 1 \end{pmatrix} u(t), \\ x(0) = x_0 \text{ unknown,} \\ y = \begin{pmatrix} 0 & 1 \end{pmatrix} x, \end{cases}$$

with  $x = (z \ \dot{z})^\top$  and  $u(t) = 5 + e^{-100t} \sin(\frac{3}{4}t)$ .

## A TEXTBOOK EXAMPLE

Consider the second order system (in matricial form)

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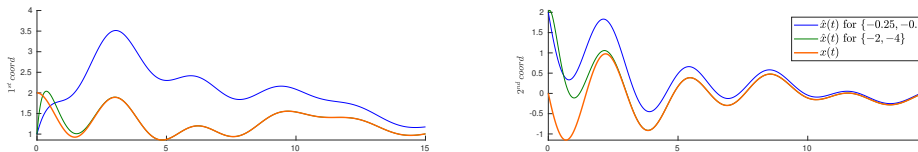
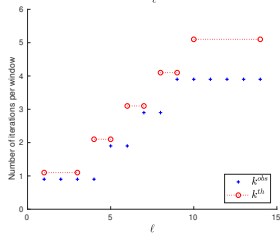
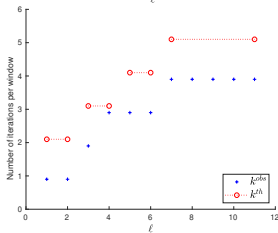
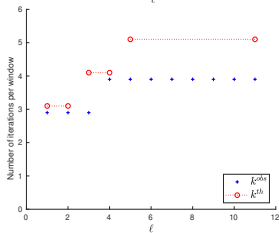
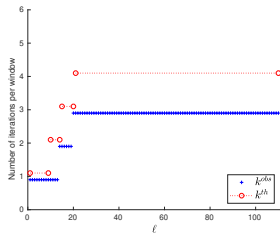
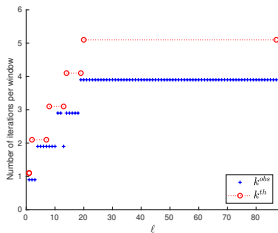
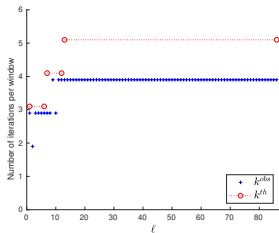


Figure:  $\eta = 0.1$ ,  $\omega = 2$



(a)  $\tilde{\gamma} = 10^{-3}$

(b)  $\tilde{\gamma} = 1$ .

(c)  $\tilde{\gamma} = 10^3$ .

**Figure:** Comparison between  $k^{th}$  and  $k^{obs}$ , for  $N = 16$  and  $\delta t = \frac{\Delta T}{2^5}$ . The eigenvalues of  $A - LC$  are  $\{-0.25, -0.5\}$  (top) and  $\{-2, -4\}$  (bottom).

We define the efficiency of the algorithm as

$$E = \frac{\tau_s}{N\tau_p}$$

where  $\tau_s$ ,  $\tau_p$  are the CPU times required to reach a given tolerance Tol by using a sequential and parallel solver, respectively; and  $N$  represents the number of available processors.

## THEOREM

The efficiency of the *Diamond Strategy* satisfies

$$E \leq \frac{\ell^* \tau_{\Delta T}^{\mathcal{F}}}{\tau_{\Delta T}^{\mathcal{F}} + N \tau_{\Delta T}^{\mathcal{G}}} \left( \sum_{\ell=0}^{\ell_{\parallel}^* - 1} k_{\ell} \right)^{-1},$$

where  $\tau_{\Delta T}^{\mathcal{F}}, \tau_{\Delta T}^{\mathcal{G}}$  corresponds to the computational time associated to one solving of (2) on a subinterval of length  $\Delta T$ , for  $\mathcal{F}$  and  $\mathcal{G}$  respectively; and  $\ell^*, \ell_{\parallel}$  denotes the number of windows required to reach a given tolerance Tol by using a sequential or parallel solver.

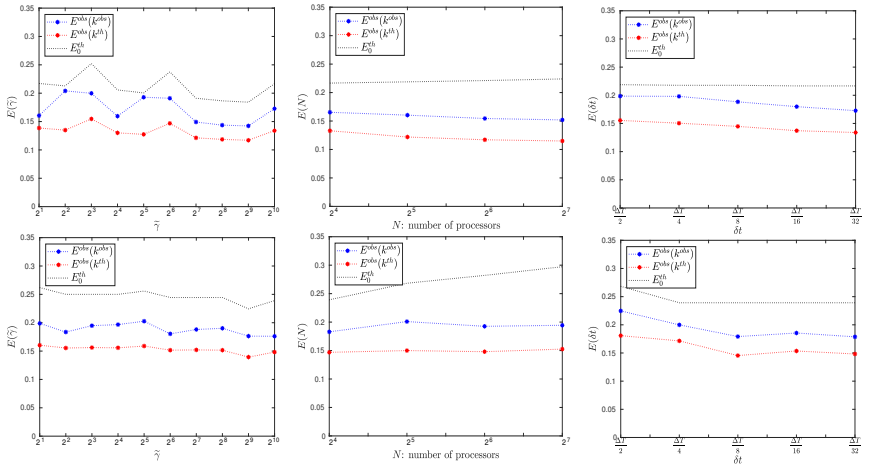


## THEOREM

The efficiency of the *Diamond Strategy* satisfies

$$E \leq E_0^{th} = \ell^* \left( \sum_{\ell=0}^{\ell_{\parallel}^* - 1} k_{\ell} \right)^{-1},$$

where  $\tau_{\Delta T}^{\mathcal{F}}, \tau_{\Delta T}^{\mathcal{G}}$  corresponds to the computational time associated to one solving of (2) on a subinterval of length  $\Delta T$ , for  $\mathcal{F}$  and  $\mathcal{G}$  respectively; and  $\ell^*, \ell_{\parallel}$  denotes the number of windows required to reach a given tolerance Tol by using a sequential or parallel solver.



(a)  $E(\tilde{\gamma})$ , for  $N = 16$  and  $\delta t = \frac{\Delta T}{2^5}$ .

(b)  $E(N)$ , for  $\delta t = \frac{\Delta T}{2^5}$  and  $\tilde{\gamma} = 2^{10}$ .

(c)  $E(\delta t)$ , for  $N = 16$  and  $\tilde{\gamma} = 2^{10}$ .

**Figure:** Comparison between  $E^{obs}(k^{obs})$ ,  $E^{obs}(k^{th})$  and  $E_0^{th}$ . The eigenvalues of  $A - LC$  are  $\{-0.25, -0.5\}$  (top) and  $\{-2, -4\}$  (bottom).

- ▶ Extension to nonlinear observers, Kalman filters.
- ▶ Use of other time-parallelization algorithms (e.g. ParaExp).
- ▶ Application to space-time problems.

Thank you for your attention !